

Lecture # 7 : Periodic Task Scheduling (DM Schedulability Analysis)

$\tau = \{\tau_1, \tau_2, \dots, \tau_n\}$ Preemption

Each $\tau_i = (C_i, T_i, D_i, \phi_i)$

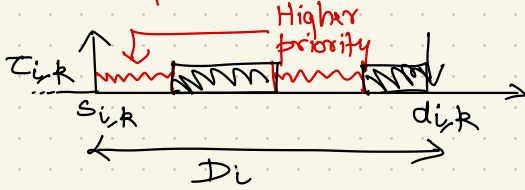
Assume: $D_1 < D_2 < \dots < D_n$

DM: $\text{prio}(\tau_1) > \text{prio}(\tau_2) > \dots$

Goal: DM schedules every task τ_i without any deadline misses

$\forall i, k$: $\tau_{i,k}$ does not miss its deadline.

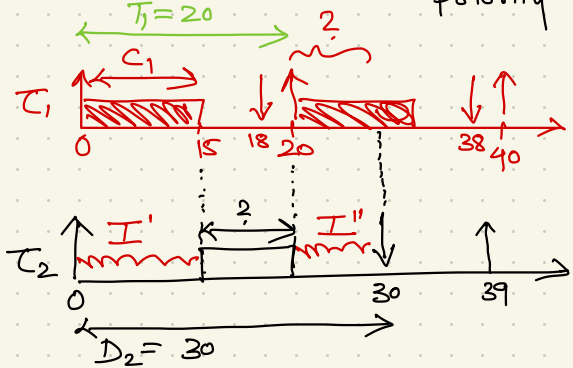
"Interference" I



$$D_i - I \geq C_i$$

$$C_i + I \leq D_i$$

τ_i	T	D	C	
τ_1	20	18	15	$D_1 < D_2$
τ_2	39	30	5	\downarrow
τ_3	100	90	8	Higher priority.



$$C_2 + I \leq D_2$$

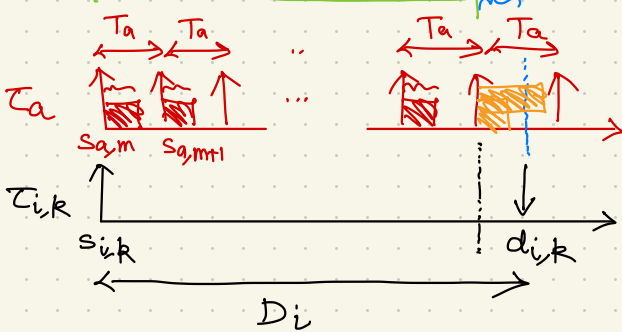
$$I' = C_1 = 15$$

$$I'' = D_2 - T_1 = 30 - 20 = 10$$

$$C_2 + I' + I'' \leq D_2$$

$$5 + 15 + 10 \leq 30? \text{ YES}$$

"FULL" $F = \lfloor D_i / T_a \rfloor$



$$C_i + I_{i,k}^a \leq D_i$$

$$\underline{I_{i,k}^{a, \text{red}}}$$

$$\underline{I_{i,k}^{a, \text{orange}}}$$

$$\underbrace{C_a \quad D_i - FT_a}_{\min(D_i - FT_a, C_a)}$$

$$C_i + I_{i,k}^{a, \text{red}} + I_{i,k}^{a, \text{orange}} \leq D_i$$

$$\forall i, k : C_i + \sum_{\substack{\text{all} \\ \text{higher} \\ \text{priority} \\ \text{tasks} \\ T_a}} (I_{i,k}^{a, \text{red}} + I_{i,k}^{a, \text{orange}}) \leq D_i$$

$0 \rightarrow \infty$

$$\forall i : C_i + \max_{\forall k} \sum_{\forall a, i} \underline{I_{i,k}^{a, \text{red}}} + \underline{I_{i,k}^{a, \text{orange}}} \leq D_i$$

$$\forall i : C_i + \sum_{\forall a < i} \max_{\forall k} \left\lfloor \frac{D_i}{T_a} \right\rfloor C_a + \min(D_i - \left\lfloor \frac{D_i}{T_a} \right\rfloor T_a, C_a) \leq D_i$$

\downarrow independent of k
 redundant

$$\forall i : C_i + \sum_{\forall a < i} \left(\left\lfloor \frac{D_i}{T_a} \right\rfloor C_a + \min(D_i - \left\lfloor \frac{D_i}{T_a} \right\rfloor T_a, C_a) \right) \leq D_i$$

$$C_3 + \sum_{a < 3} \left(\left\lfloor \frac{D_3}{T_a} \right\rfloor C_a + \min(D_3 - \left\lfloor \frac{D_3}{T_a} \right\rfloor T_a, C_a) \right) \leq D_3$$

$$C_3 + \overbrace{\left[\left\lfloor \frac{D_3}{T_1} \right\rfloor C_1 + \min(D_3 - \left\lfloor \frac{D_3}{T_1} \right\rfloor T_1, C_1) \right]}^{a=1} + \underbrace{\left[\left\lfloor \frac{D_3}{T_2} \right\rfloor C_2 + \min(D_3 - \left\lfloor \frac{D_3}{T_2} \right\rfloor T_2, C_2) \right]}_{a=2} \leq D_3$$

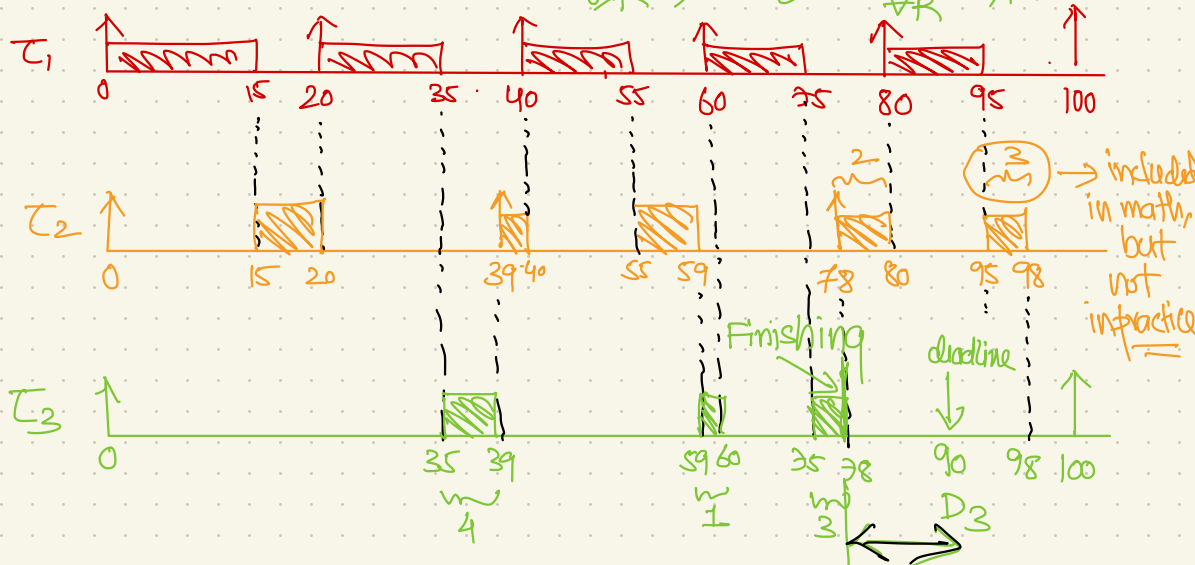
$$8 + 4 * 15 + \min(90 - 4 * 20, 15) = 8 + 60 + 10 = 93 \leq 90$$

$$+ 2 * 5 + \min(90 - 7 * 8, 5) = + 10 + 5$$

NO!
 \downarrow

Finishing time - starting time = R

$R_{ik} \text{ ? } R_i = \max_{k \neq i} R_{ik}$



Old test : Sufficient but not necessary

$$C_i + I \leq D_i \Rightarrow C_i + \underline{I} \leq R_i$$

↳ also depends on R_i

Response-Time Analysis

RTA

HW: Sec 4.5.1 & 4.5.2